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<p>(21) International Application Number: PCT/CA92/00152</p> <p>(22) International Filing Date: 14 April 1992 (14.04.92)</p> <p>(30) Priority data: 684,775 15 April 1991 (15.04.91) US</p> <p>(71) Applicant: CANAI COMPUTER AND NETWORK ARCHITECTURE INC. [CA/CA]; 59 Iber Road, Stittsville, Ontario K2S 1E7 (CA).</p> <p>(72) Inventors: DELANEY, David, M. ; 142 Waverley Street, Ottawa, Ontario K2P 0V4 (CA). COTTREAU, Peter, M., K. ; 4-433 Besserer Street, Ottawa, Ontario K1N 6B9 (CA). CHEFURKA, Paul, V. ; 1020 Barwell Avenue, Ottawa, Ontario K2B 8H5 (CA). GLAVIN, Lisanne, M. ; 173 Walden Drive, Kanata, Ontario K2K 2K8 (CA).</p>			<p>(74) Agent: PASCAL, Edward, E.; P.O. Box 11121, Station H, Ottawa, Ontario K2H 7T8 (CA).</p> <p>(81) Designated States: AT, AT (European patent), AU, BB, BE (European patent), BF (OAPI patent), BG, BJ (OAPI patent), BR, CF (OAPI patent), CG (OAPI patent), CH, CH (European patent), CI (OAPI patent), CM (OAPI patent), CS, DE, DE (European patent), DK, DK (European patent), ES, ES (European patent), FI, FR (European patent), GA (OAPI patent), GB, GB (European patent), GN (OAPI patent), GR (European patent), HU, IT (European patent), JP, KP, KR, LK, LU, LU (European patent), MC (European patent), MG, ML (OAPI patent), MN, MR (OAPI patent), MW, NL, NL (European patent), NO, PL, RO, RU, SD, SE, SE (European patent), SN (OAPI patent), TD (OAPI patent), TG (OAPI patent).</p> <p>Published <i>With international search report Before the expiration of the time limit for amending the claims and to be republished in the event of the receipt of amendments.</i></p>
<p>(54) Title: ADDRESS CACHING NETWORK INTERFACE CIRCUIT FOR TRANSPARENT BRIDGE</p> <p>(57) Abstract</p> <p>A data transfer system between local area networks to which stations are connected comprising a transparent bridge (10), network interface circuits (NICs) (16, 20) connected between the networks and the bridge, at least one NIC comprising cache memory apparatus (22) for maintaining temporary record of the source address of station(s) transmitting data on an associated LAN, apparatus for determining the destination address of frames received on the LAN from a given station or another bridge attached to the LAN, apparatus for transferring (or allowing the transfer of) the frames to the bridge for further transmission to another NIC in the event the destination address is not one of the addresses in the cache memory.</p>			

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ADDRESS CACHING NETWORK INTERFACE CIRCUIT
FOR TRANSPARENT BRIDGE

Field of the Invention

5 This invention relates to data transmission networks, and in particular to a system for transferring data between local area networks (LANs).

Background to the Invention

10 A local area network carries data between stations connected to the network. Examples of local area networks are, for example, the Arcnet LAN, the Ethernet LAN, the IBM Token Passing LAN, etc. In order to enable the communication of data from one station on one LAN to another station on another LAN, a bridge is 15 connected between the LANs.

20 A simplified LAN configuration is shown in Figure 1. The bridge 10 forms a point at which data traffic (frames) is transferred by the bridge between the three LANs 11, 12, and 13 in order that stations connected to different LANs may communicate with each other. Each LAN is connected to the bridge by means of a network interface circuit, or NIC, 14, 15, 16. The NIC is usually, but not necessarily, a printed circuit card distinct from the other parts of the bridge, and 25 is usually connected to the LAN by means of a backplane component of the bridge or by other means, such as a connector on the main printed circuit board of the bridge.

30 Each NIC performs the functions necessary to connect the bridge to a particular type of LAN; each NIC typically includes all of the hardware necessary to connect to the specific type of LAN for which it is designed. Several LANs of different types may be connected in this way to a single bridge.

35 Although NICs are usually considered to be components of the bridge in which they are found, the term bridge is used herein to refer to the bridge exclusive of the NICs.

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Description of the Prior Art

In the prior art, a NIC attached to a transparent bridge performs promiscuous reception. A NIC performs promiscuous reception by transferring to its bridge all frames it receives from the LAN.

Promiscuous reception is appropriate only in a bridge, or in a special station which has the responsibility of gathering statistics about the operation of the LAN. A NIC operating in an ordinary station transfers to the station only those frames which are addressed to the station.

A bridge processes data frames received from its NICs by utilizing addresses in the structure of the frames. Addresses in IEEE LANs (and some other LANs, such as FDDI) conform to the IEEE standard for LAN addresses described in IEEE standard 802.1: each frame has a destination address field and a source address field (see Figure 2). Both destination and source addresses have the same standard format.

Each NIC attached to a LAN must have an individual IEEE address which is unique within the LAN to which the NIC is attached, and which is also unique within the extended LAN formed by that LAN and other LANs to which it is connected, either directly or indirectly, by bridges.

The destination address field may contain either an individual address, or a group address. The first bit of an address determines whether the address is an individual address or a group address. The first bit of an individual address is 0; the first bit of a group address is 1.

A station places the unique individual address of its NIC into the source address field of each frame it originates. Only individual addresses may appear in the source address field. In the destination address field, the station places either the unique individual address of a single NIC, or a group address.

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In normal operation, a NIC attached to a station will pass a frame from the LAN to the station only if the destination address field of the frame contains either the individual address of the NIC, or a group address to which the station belongs. The use of an individual address in the destination address field of a frame results in only a single station receiving the frame; the use of a group address results in potentially many stations receiving the frame. Some prior art NICs have a feature for filtering group addresses by hashing addresses found in the destination field into a bit location in a hash register in the NIC. The bits of the hash register of the NIC can be set individually by the station software. The station software sets bits corresponding to group addresses which it should receive. The NIC passes a frame bearing a group address to the station if the bit of the hash register corresponding to the group address is set.

In the prior art, a NIC attached to a bridge performs promiscuous reception, ignoring both the destination and source address fields of received frames, and transferring every frame to the bridge. A NIC attached to a station inspects the destination address field of each frame received from the LAN, and transfers a frame to the station when the destination address of the frame is identical to the address of the NIC, or when the destination address of the frame is a group address.

A transparent bridge builds and maintains an internal routing table. Each entry of the routing table has two components: an individual address designating a station somewhere on the extended LAN, and a number designating an outgoing NIC for the address. The outgoing NIC is one of the NICs attached to the bridge. The outgoing NIC for an address must be connected to a LAN (or extended LAN) to which the NIC bearing the address is attached.

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The NIC of a bridge through which the bridge receives a frame is the source NIC of the frame.

A transparent bridge builds and maintains its routing table by inspecting the source address fields of frames received from its NICs. (Hence the alternative name learning bridge.) If a frame bearing a given address in its source address field is received via a given NIC, then the bridge may transmit other frames destined for the given address through that NIC. Thus the following procedure is performed on each frame received.

The bridge determines whether the source address of the received frame is in the routing table. If the address is not in the table, the bridge creates a routing table entry containing the source address and the number of the source NIC.

If the destination address of a received frame is a group address, the bridge transmits the frame through all of its NICs except the source NIC.

If the destination address of a received frame is an individual address, the bridge routes the frame as follows:

1. The bridge determines whether the destination address is in the routing table.
2. If the destination address of the frame is not present in the table, the bridge transmits the frame through all of its NICs except the source NIC.
3. Otherwise, if the outgoing NIC for the destination address (the NIC specified in the routing table entry for the destination address) is the same as the source NIC, the bridge discards the frame.
4. Otherwise, the source and outgoing NICs are different, and the bridge transmits the frame through the outgoing NIC.

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5 The IEEE transparent bridging standard, IEEE 802.1d, specifies procedures (the so-called spanning tree algorithm) by which the bridges of an extended LAN manage their routing tables so that there will be only a single path which frames may follow between any two points of an extended LAN. This property ensures that a given source address will be observed in frames received by only one of the NICs of a bridge.

10 In order to permit physical reconfiguration of the extended LAN, routing table entries are aged. A given entry is permitted to remain in the table only a few minutes before it is deleted. IEEE 802.1d recommends five minutes as the default lifetime permitted for a routing table entry.

15 Summary of the Invention

20 The present invention enhances the performance of transparent bridges by introducing a new kind of NIC. In a prefiltering function, this NIC discards frames that the bridge would otherwise discard before the frames reach the bridge.

25 The bridge is spared the requirement to copy every frame seen by the NIC from the NIC into the bridge, and the requirement to perform two address lookups for the frame, one for the destination address and one for the source address. This allows the bridge processor to handle more NICs than in the prior art. Significant enhancement of the bridge-NIC combination is thus obtained.

30 Instead of promiscuous reception, the NIC of the present invention performs selective reception. Bridge processing steps need not be affected by the difference between selective reception and promiscuous reception. Attaching a NIC which embodies selective reception does not require modifying the existing bridge procedures; such NICs may be fitted to existing bridges with little effort. In this case, the unmodified transparent bridge may handle an increased number of NICs without blocking, since the bridge need

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not handle data frames having both source and destination addresses related to the same LAN, in many cases. The exceptions will be described below.

A NIC with bridge-assistance hardware 5 operates as follows: The NIC has a source address filter cache (SAFC) containing up to some fixed number (e.g. 256) individual source addresses found in frames received from the LAN. As the network interface card receives each frame, it searches the SAFC for the source address of the frame and for the destination address of the frame. If the source address is not present in the SAFC, the NIC inserts the source address 10 into the SAFC and passes the frame to the bridge, ensuring that the bridge always discovers which frames are present on the LAN. If the source address is 15 present in the SAFC, the NIC discards the frame if the destination address is present in the SAFC, and otherwise passes the frame to the bridge.

When the NIC passes a frame to the bridge it 20 provides an indication of the results of this processing in a so-called maintenance word which accompanies the frame in the frame buffer. The bridge can find out by checking the value of two bits in the maintenance word whether each of the source and destination addresses was found in the SAFC of the NIC. 25 The bridge updates its routing table only when the maintenance word of a frame indicates that the NIC did not find the source address of the frame in its SAFC. After checking the indicator for the source address and possibly (infrequently) updating the routing table, the 30 bridge discards any frame for which the destination address of the frame was found in the SAFC of the NIC. Frames which remain after this check may have to be forwarded or discarded, but only one further table 35 lookup is required to determine whether the frame is to be forwarded or discarded. The bridge looks up the destination address in the routing table, if the address is not found, the frame is flooded to all LANs

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connected to the bridge except the LAN on which the frame was received. If the destination address is found in the routing table, the bridge forwards the frame to the LAN specified in the entry in the routing
5 table, unless the specified LAN is the one on which the frame was received, in which case the frame is discarded.

To summarize, the presence of the bridge-assistance feature in all NICs connected to a bridge
10 reduces the load on the bridge CPU in two ways, as follows:

- (1) The NICs autonomously discard many frames, preventing the CPU from expending any time whatsoever on such
15 frames.
- (2) The NICs provide the CPU with an efficient indication of whether the source address needs to be looked up by the CPU. In situations where forwarding is frequent, the source address will almost always have been found in the SAFC of the NIC, and no source address lookup will be required by the CPU. It
20 is precisely when forwarding is most frequent that the bridge-assistance feature is most effective and most needed to reduce the CPU resource
25 required for the forwarding task.

An embodiment of the present invention is a data transfer system between local area networks to which stations are connected comprising a transparent bridge, NICs connected between the networks and the bridge, at least one NIC comprising: a cache memory for maintaining temporary record of the source address of a station transmitting data on the associated LAN,
30 apparatus for determining the destination address of data received on the LAN, an apparatus for transferring
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or allowing the transfer of the data to the bridge for further transmission to another NIC in the event the destination address is not one of the addresses in the cache memory.

5 In accordance with a particular form of the invention, it is determined whether the source address of a frame of data is stored in the cache memory, and the data is discarded in the event the destination address is the same as a source address stored in the cache memory. If the source address of a data frame is not stored in the cache memory, it is stored and the frame is passed to the bridge. For subsequent frames the source address is determined, and the system operates as above.

10 Brief Introduction to the Drawings

A better understanding of the invention will be obtained by reference to the detailed description below, in conjunction with the following drawings, in which:

20 Figure 1 is a block diagram of a LAN data transfer system in accordance with the prior art,

Figure 2 is a diagram of the configuration of a frame of data,

25 Figure 3 is a truth table relating to the operation of the invention,

Figure 4 is a block diagram of the present invention, and

Figure 5 is a logic diagram of a portion of the present invention.

30 Description of the Preferred Embodiments

In accordance with the present invention, a NIC possesses a source address cache memory. Briefly, the NIC places the source address of a frame received from the LAN into the source address cache and passes the frame to the bridge, if the address is not already present in the cache, and otherwise discards a frame received from the LAN if the destination address is found in the source address cache.

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The addresses are located in the part of the frame which is received first. The relevant aspects of the format of frames of IEEE LANS and FDDI are displayed in Figure 2. It may be seen that each frame 5 is comprised of a frame destination address, a frame source address, the data bits to be transferred, and a frame check sequence (FCS).

Figure 4 is a block diagram of a structure 10 for implementing the present invention. The structure is similar to that of Figure 1, except that one of the NICs 14 is replaced by a NIC 20. The NIC 20 contains a cache memory 22, as will be described below.

The bridge 10, the same as in the prior art, 15 contains in a memory a routing table 24 which contains a source address column 26 and a corresponding source NIC identification number column 28. Each source address has an entry in column 26 opposite a source NIC identification number in column 28.

The source addresss cache memory 22 is 20 preferably a content addressable memory such as type AM99C10A produced by Advanced Micro Devices (AMD). The performance of the cache must be such as to permit two lookups for every frame at the maximum rate at which frames may be received from the LAN, and leave time to 25 perform cache management activities and other processing associated with the frame. For example, approximately fifty thousand lookups per second are required for connection to an Ethernet LAN, and five hundred thousand lookups per second are required for 30 connection to a FDDI LAN.

As soon as the destination and source 35 addresses of each frame arrive in the NIC from the LAN, the NIC starts to perform an address processing procedure, described later. Address processing includes looking up the source and destination addresses in the source address cache, and deciding on the disposition of the frame. The actions consequent upon address processing (inserting the source address

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in the cache, forwarding the frame to the bridge, etc) should not be performed until the Frame Check Sequence (FCS) (Cycle Redundancy Calculation) at the end of the frame has been received and validated by the NIC, since
5 the addresses may have been corrupted. If the frame check calculation indicates that the frame has been corrupted, the frame is discarded and none of the pending actions consequent upon the address processing are performed. If the frame check calculation
10 indicates that the frame has not been corrupted, the pending actions consequent upon address processing are then performed. Address processing is preferably concurrent with the reception of the portion of the frame subsequent to the addresses, ensuring that
15 address processing will be complete, and that consequent actions will have been determined, and will be ready for execution, by the time the result of the FCS calculation becomes available at the end of the frame.

20 In order to implement the above, as well as to implement the functional steps described below, reference is made to the block diagram shown in Figure 5. A data frame carried on LAN 11 appears at the input of a frame buffer 30, a frame address discriminator 31, and an FCS checker 35, connected at the input of the
25 NIC 20 which is connected to the LAN 11. Buffer 30 is of sufficient size to hold an entire frame of data. Frame address discriminator 31 passes the destination address and the source address of each received packet to the address cache 22 via cache address bus 32 which is shared between the frame address discriminator and the NIC controller 100.

30 The preferred embodiment utilizes a content addressable memory (CAM), such as type AM99C10A, made by Advanced Micro Devices, Inc., as the source address cache 22. A typical CAM has two modes of access, the lookup mode and the direct access mode. The CAM is arranged as a linear sequence of memory locations. In
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the lookup mode, the CAM is presented with a search argument and responds with a signal specifying whether the search argument is present in one of the memory locations of the CAM. In the direct access mode, 5 individual locations of the CAM may be read, written, or made empty.

The NIC uses the lookup mode to determine whether the destination and source addresses of the frame are present in the source address cache. The 10 status of the address lookup operation is communicated to the NIC controller 100 via an address match status signal on path 103. The NIC uses the direct access mode to conditionally insert source addresses into the source address cache. The NIC controller signals the 15 source address cache when the source address of the frame is to be inserted into the cache via an insert source address control signal. The NIC controller maintains a pointer to the sequence of memory locations of the source address cache. When the NIC inserts a 20 source address into the source address cache it does so at the location specified by its pointer into the cache which the NIC controller communicates to the source address cache via the cache address bus 32. The NIC controller then increments the pointer so that it 25 points at the next location of the CAM, resetting the pointer so that it points to the first location of the CAM if the incrementation results in pointing farther than the last location of the CAM. Newly inserted source addresses eventually overwrite old source 30 addresses.

As the destination and source address lookups are proceeding, the frame is written into the frame buffer 30. The frame data is also processed by the FCS checker 35. The last bits in the frame are the frame 35 check sequence which are also stored in the buffer 30, and processed by the FCS checker 35. The status of the FCS verification is communicated to the NIC control logic via the FCS status signal path 102.

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Some LANs, such as FDDI, permit the reception by a NIC of frames transmitted to the LAN by the same NIC. If the LAN permits frames transmitted by a station and addressed to itself to be received by the same station, then frames having the same address in both the source and destination address fields must not be discarded. In a NIC for such a LAN, the NIC determines the value of the predicate "source address equals destination address equals self address" (SDS).
5 The NIC uses the resource of the self address comparator 106 to generate the SDS predicate. The status of the SDS predicate is communicated to the NIC control logic via the SDS status signal 107.
10

Once the frame reception is complete then, 15 assuming that the FCS checks correctly, the NIC control logic manipulates the resource of the NIC based on status signals on paths 102, 103 and 107 according to the truth table given in Figure 3.

In the process, with reference to the truth 20 table of Figure 3 carried out by the NIC, SDS may be considered to be FALSE always, if the type of the LAN is such that frames sent by a station and addressed to itself cannot be received by the station. Note that SDS cannot be TRUE for a frame bearing a group address 25 in its destination address field.

For a given frame, the predicate "source 30 address in cache" (SC) is TRUE if and only if the source address of the frame is present in the source address cache. The NIC computes the value of SC by looking up the source address of the frame in the source address cache of the NIC.

For a given frame, the predicate "destination 35 address in cache" (DC) is TRUE if and only if the destination address of the frame is present in the source address cache. The NIC computes the value of DC by looking up the destination address of the frame in the source address cache of the NIC.

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The action by which the electronic logic of the NIC forwards a frame to the bridge is denoted FF.

The action by which the electronic logic of the NIC discards a frame so that the bridge does not become aware of its reception is denoted DF.

The action by which the electronic logic of the NIC inserts the source address of a frame into the source address cache is denoted ISC.

With the above definitions, the table of 10 Figure 3 specifies the address processing procedure, as follows:

1. If the address of the NIC is the same as both the destination and source addresses of a frame received from the LAN, then forward the frame to the bridge.
2. Otherwise, if the source address of the frame is present in the source address cache, then forward the frame to the bridge if the destination address of the frame is NOT in the source address cache, otherwise discard the frame.
3. Otherwise the source address is NOT in the source address cache; the source address should be inserted into the source address cache, and the frame should be forwarded to the bridge.

Forwarding a frame whenever the source address is not present in the source address cache 30 ensures that the bridge learns the location of all source addresses.

The above is correct provided that group addresses are never found in the source address cache. The presence of a group address in the source address cache would result in frames destined for that group address being discarded by the NIC, which would be wrong, since the bridge must receive all such frames. Before address processing is performed, the NIC should

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5 discard all frames which have a 1 in the first bit of
the source address field, thus ensuring the required
condition even in the presence of errors or source
routed frames, which use the first bit of the source
address field as a flag to indicate that a route has
been inserted in the frame following the source
address.

10 The maximum number of source addresses that
the source address cache may hold can be substantially
smaller than the number of active station on the
extended LAN. The bursty nature of the traffic created
by or for any particular station ensures that even a
relatively small cache, say 256 entries, should be
adequate for a very large LAN of several thousand
15 stations.

20 In order to meet the requirement of the
IEEE 802.1d standard that individual addresses should
reside in the routing table of the bridge for only a
limited time, the lifetimes of addresses in the source
address cache of the NIC must also be limited. The NIC
cannot rely on the overwriting of old addresses by new
ones, since too few distinct source addresses may be
received by the NIC within the lifetime permitted for
entries in the cache. The direct access mode of the
25 CAM may be used to remove entries from the source
address cache as follows:

30 The NIC or the bridge maintains a pointer (a
deletion pointer) to a location of the CAM. At regular
intervals, the location of the CAM specified by the
pointer is made empty using the direct access mode of
the CAM. The pointer is then incremented to point to
the following location of the CAM, and is reset to
point to one end of the CAM if incrementation results
in pointing outside of the CAM. The frequency of this
process is chosen so that every location of the CAM
35 will be erased within a period less than the lifetime
permitted for residence of an address in the routing
table of the bridge. There is no need to track the age

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of individual entries of the source address cache. A slight benefit can be obtained by incrementing the deletion pointer in the direction opposite to that in which the insertion pointer is incremented.

5

Variations or other embodiments of the invention may now be designed, using the principles described above. All those which fall within the bounds of the claims are considered to be part of the present invention.

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WE CLAIM:

1. A data transfer system between local area networks (LANs) to which stations are connected comprising:

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(a) a transparent bridge,

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(b) network interface circuits (NICs), one connected between each LAN and said bridge,

15

(c) at least one NIC comprising

20

(i) cache memory means for maintaining temporary record of the source address of stations transmitting data on the associated LAN,

25

(ii) means for determining the destination address of data received on the LAN from said station,

30

(iii) means for determining the source address of data received on the LAN from the station,

(iv) means for determining whether the source address in the data is stored in the cache memory, and for discarding the data in the event the destination address is the same as a source address stored in the cache memory, and

(v) means for transferring said data to the bridge for further transmission to another NIC in the event the destination address is not the same as a source address stored in the cache memory.

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2. A system as defined in claim 1, further including means for storing the source address in the cache memory in the event the source address of data received from the LAN is not stored in the cache memory.

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3. A system as defined in claim 2, further including means for maintaining particular addresses in the cache memory for a limited period of time.

5

4. A method of transferring data between a first LAN and a transparent bridge comprising receiving data carried on the LAN, prefILTERING the data to remove frames of data destined for stations connected to said first LAN, and transferring or allowing transfer of data destined for stations connected to another LAN to the transparent bridge for further transfer to said another LAN, thereby reducing the amount of data to be dealt with by the transparent bridge.

10

15

5. A method as defined in claim 4 including the step of temporarily storing source addresses of data received from said first LAN in a cache memory, and in the prefILTERing step comparing the destination address of data received from said first LAN with said stored source addresses to determine whether the data should or should not be transferred or allowed to be transferred.

20

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30

6. A method as defined in claim 5, including the step of storing the source address in the cache memory in the event the source address of the data received from said first LAN is not stored in the cache memory.

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7. A method as defined in claim 6, including the step of clearing the cache memory after source addresses have been stored for a predetermined period of time.

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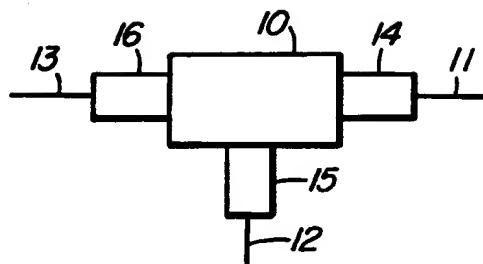


FIG. 1

FIG. 2

DESTINATION ADDRESS	SOURCE ADDRESS	DATA	FCS
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SDS	SC	DC	ACTION
T	X	X	FF
F	T	F	FF
F	T	T	DF
F	F	X	ISC, FF

T = TRUE
F = FALSE
X = DON'T CARE

FIG. 3

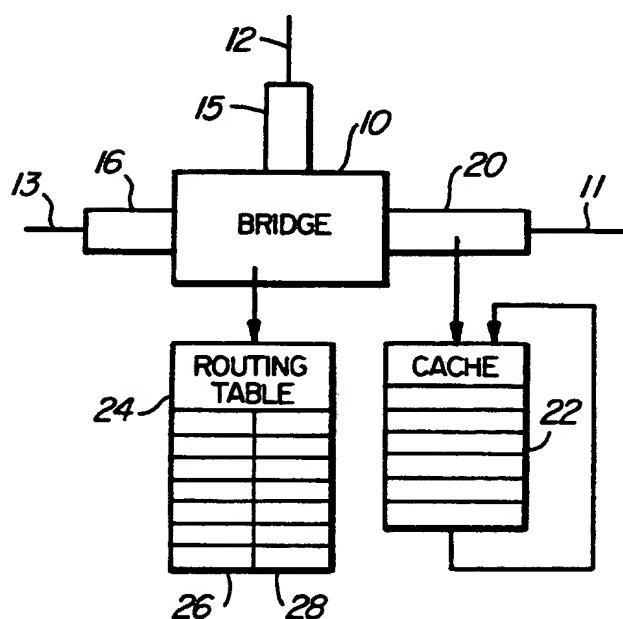


FIG. 4

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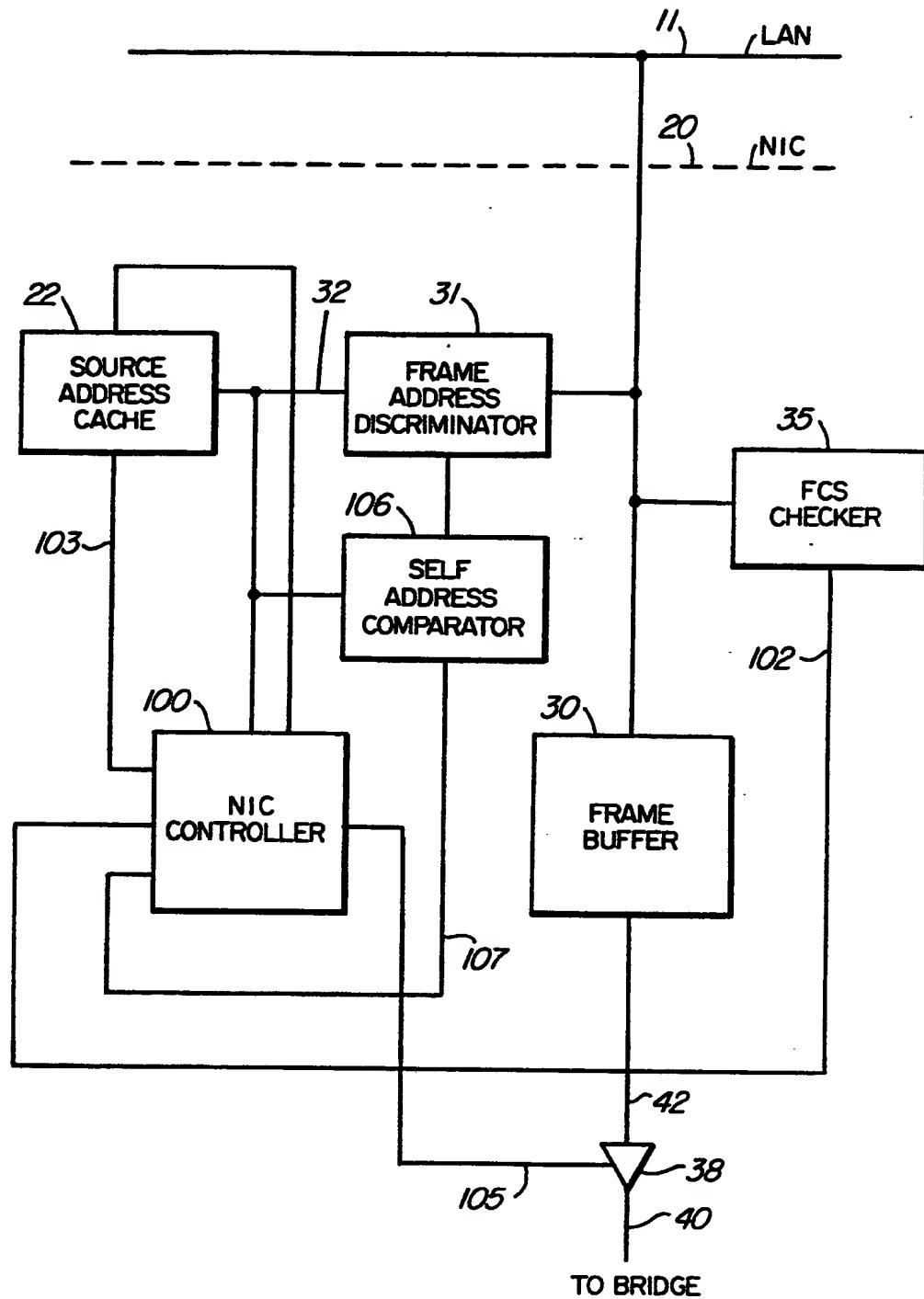


FIG. 5

INTERNATIONAL SEARCH REPORT

International Application No

PCT/CA 92/00152

I. CLASSIFICATION OF SUBJECT MATTER (If several classification symbols apply, indicate all)⁶

According to International Patent Classification (IPC) or to both National Classification and IPC

Int.C1. 5 H04L12/46

II. FIELDS SEARCHED

Minimum Documentation Searched⁷

Classification System	Classification Symbols
Int.C1. 5	H04L

Documentation Searched other than Minimum Documentation
to the Extent that such Documents are Included in the Fields Searched⁸III. DOCUMENTS CONSIDERED TO BE RELEVANT⁹

Category ¹⁰	Citation of Document, ¹¹ with indication, where appropriate, of the relevant passages ¹²	Relevant to Claim No. ¹³
X	US,A,4 922 503 (INFOTRON SYSTEMS) 1 May 1990 see column 2, line 61 - column 3, line 16 see column 6, line 61 - column 9, line 9 ---	4 1-3,5-7
Y	EP,A,0 156 542 (ICL) 2 October 1985 see page 2, line 15 - page 4, line 26 see page 11, line 21 - line 22 ---	1-3,5-7
A	EP,A,0 222 584 (UNIVERSITY OF SALFORD) 20 May 1987 see column 1, line 20 - line 33 ---	2,6
A	IBM TECHNICAL DISCLOSURE BULLETIN. vol. 33, no. 7 December 1990, NEW YORK US pages 63 - 68; 'Bridge for Interconnecting Local Area Networks Having Different Architectures' see page 64, line 26 - line 40 ---	3,7

¹⁰ Special categories of cited documents :¹⁰

- "A" document defining the general state of the art which is not considered to be of particular relevance
- "E" earlier document but published on or after the international filing date
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IV. CERTIFICATION

Date of the Actual Completion of the International Search

Date of Mailing of this International Search Report

1. 09 SEPTEMBER 1992

01.10.92

International Searching Authority

Signature of Authorized Officer

EUROPEAN PATENT OFFICE

MESSELKEN M.

**ANNEX TO THE INTERNATIONAL SEARCH REPORT
ON INTERNATIONAL PATENT APPLICATION NO.**

CA 9200152
SA 58363

This annex lists the patent family members relating to the patent documents cited in the above-mentioned international search report.
The members are as contained in the European Patent Office EDP file on
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09/09/92

Patent document cited in search report	Publication date	Patent family member(s)			Publication date
US-A-4922503	01-05-90	None			
EP-A-0156542	02-10-85	AU-B-	569380	28-01-88	
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